

Bi-criteria Scheduling on Parallel Machines Under Fuzzy Processing Time

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Abstract: Job scheduling is concerned with the optimal allocation of scarce resources with objective of optimising one or several criteria. Job scheduling has been a fruitful area of research for many decades in which scheduling resolve both allocation of machines and order of processing. If the jobs are scheduled properly, not only the time is saved but also efficiency of system is increased. The parallel machine scheduling problem is widely studied optimization problem in which every machine has same work function and a job can be processed by any of available machines. Optimising dual performance measures on parallel machines in fuzzy environment is fairly an open area of research. In real life situations, the processing times of jobs are not always exact due to incomplete knowledge or an uncertain environment which implies the existence of various external sources and types of uncertainty. Fuzzy set theory can be used to handle uncertainty inherent in actual scheduling problems.

This paper pertains to a bi-criteria scheduling on parallel machines in fuzzy environment which optimizes the weighted flow time and total tardiness simultaneously. The fuzziness, vagueness or uncertainty in processing time of jobs is represented by triangular fuzzy membership function. The objective of the paper is to find the optimal sequence of jobs processing on parallel machines so as to minimize the secondary criterion of weighted flow time without violating the primary criterion of total tardiness. The bi-objective problem with total tardiness and weighted flow time as primary and secondary criteria respectively, for any number of parallel machines is NP-hard. A numerical illustration is carried out to the test efficiency of the proposed algorithm.

Keywords: *Fuzzy processing time, total tardiness, weighted flow time, due date, weighted job*

1. INTRODUCTION

Job scheduling has been a fruitful area of research for many decades in which scheduling resolve both allocation of machines and order of processing. If the jobs are scheduled properly, not only the time is saved but also efficiency of system is increased. The parallel machine scheduling problem is widely studied optimization problem in which every machine has same work function and a job can be processed by any of available machines. Optimising dual performance measures on parallel machines in fuzzy environment is fairly an open area of research. A survey of literature has revealed little work reported on the bicriteria scheduling problems on parallel machine in fuzzy environment. Chen and Bulfin (1989) have examined single machine scheduling problems when all jobs have identical processing time. Cenna and Tabucanon (1991) dealt with bicriteria scheduling with parallel identical machine minimizing the maximum tardiness. Chen and Bulfin (1994) studied scheduling on a single machine to minimize the maximum tardiness (maximum latness) and number of tardy jobs (late jobs). Parkash (1997) studied the bi-criteria scheduling problems on parallel machines. Fiala (1997) formulated a parallel machine scheduling problem with two criteria to minimize the makespan and to minimize the number of job preemption. Sarin and Hariharan (2000) considered the bicriteria problem of scheduling ' n ' jobs on two machines to minimize the primary criterion of maximum tardiness and secondary criterion of number of tardy jobs. Sarin and Parkash (2004) consider the problem of scheduling jobs on parallel identical machines so as to minimize primary and secondary criteria. Anghinolfi and Paolucci (2007) studied total tardiness scheduling problems on parallel machines. Huo et al (2007) studied bicriteria scheduling problems involving number of tardy jobs and maximum weighted tardiness. Gupta et al (2012) developed an algorithm for the bi-objective problem which optimizes the number of tardy jobs without violating the value of maximum tardiness with uncertain processing time. Sharma et al (2013) studied the bi-objective problem with total tardiness and number of tardy jobs as primary and secondary criteria respectively for any number of parallel machines in fuzzy environment. Sharma et al (2013a) developed an algorithm to schedule jobs on parallel identical machines so as to minimize the secondary criterion of weighted flow time without violating the primary criterion of maximum tardiness in fuzzy environment.

The rest of paper is organized as follows: In section 2 problems is formulated. Section 3 describes the role of fuzzy in scheduling. Section 4 deals with the theorems derived for optimising the bicriteria problem. Section 5 defines the algorithm proposed to find the optimal sequence for bicriteria problem Weighted flow time/Total tardiness. In section 6, numerical illustrations are carried out to test the efficiency of the proposed algorithm. The paper is concluded in section 8 followed by the references.

2. PROBLEM FORMULATION

The following assumptions are made before proceeding with the mathematical formulation in developing the algorithm for bicriteria problem on parallel machines.

- Jobs are available at time zero
- Jobs are independent of each other
- No pre-emption is allowed
- Machines are identical in all respects and are available all the time
- No machine can handle more than one job at a time.

The following notations will be used all the way through the present paper.

- | | |
|---|---|
| • i : i^{th} job, $i = 1, 2, 3, \dots, n$ | • n : number of jobs to be scheduled |
| • d_i : due date of the i^{th} job | • m : number of machines |
| • c_i : completion time of i^{th} job | • t_a : completion time for job 'a' |
| • w_i : weight of i^{th} job | • k : machine on which i^{th} job is assigned at the j^{th} position |
| • T_i : tardiness of the i^{th} job | • WFT : weighted flow time of jobs |
| • T_{max} : maximum tardiness | • X_{ijk} : 1; if job i is located at position j on k^{th} machine and 0; otherwise |
| • j : location of i^{th} job on machine k , where $j = 1, 2, 3, \dots, n$ | |

Before formulating the bicriteria problem, the mathematical formulation for the single criterion is represented first as given by Parkash (1997). These are as discussed in section 2.1:

2.1 Criterion: Total Tardiness

Tardiness is given by $T_i = \max(0, c_i - d_i)$, where c_i and d_i are the completion time and due date of job i , respectively. The formulation is as follows:

$$\min Z = \sum_{i=1}^n T_i$$

Subject to constraints:

$$\sum_{j=1}^n \sum_{k=1}^m X_{ijk} = 1 \quad \forall i \quad \text{--- (1)}$$

$$\sum_{i=1}^n X_{ijk} \leq 1 \quad \forall j, k \quad \text{--- (2)}$$

$$X_{ijk} \text{ is binary} \quad \forall i, j, k \quad \text{--- (3)}$$

$$T_i \geq c_i - d_i \quad \forall i \quad \text{--- (4)}$$

along with non-negativity constraints.

2.2 Criterion: Weighted Flow time

The formulation to minimize the weighted flow time (WFT) is as follows:

$$\min Z = \sum_{i=1}^n \sum_{j=1}^m \sum_{k=1}^{m!} w_i \cdot X_{ijk}$$

Subject to: constraints set (1), (2) and (3) respectively, along with non-negativity constraints.

The formulation of the bicriteria problems is similar to that of single criteria problems but with some additional constraints requiring that the optimal value of the primary objective function is not violated. The two parts of the bicriteria problem formulation are as follows:

- **Primary objective function**
Subject to: Primary problem constraint
- **Secondary objective function**

Subject to:

- a. secondary problem constraint
- b. primary objective function value constraint
- c. primary problem constrain

Here, the problem is divided in two steps: one is primary in which total tardiness of jobs is minimized and in secondary step: the weighted flow time of jobs is minimised under the objective function value of primary criterion.

3. ROLE OF FUZZY

Fuzzy set theory has been used to model systems that are hard to define precisely. Zadeh (1965) stated that most of the early interest in fuzzy set theory pertained to representing uncertainty in human cognitive system. Uncertainty can be thought of in an epistemological sense as being the inverse of information. Information about a particular problem may be incomplete, imprecise, fragmentary, unreliable, vague or deficient in some other way. Fuzzy set theory is now applied to problems in engineering, business, medical and related health sciences and in natural sciences. A large number of deterministic scheduling algorithms have been proposed in last decades to deal with scheduling problems with various objectives and constraints. In real life situations, decisions to be made are often constrained by specific requirements. The decision making process gets increasingly more complicated with increments in the number of constraints. The real world is complex; complexity in the world generally arises from uncertainty. From this prospective the concept of fuzzy environment is introduced in the field of scheduling. For example, the processing times of jobs may be uncertain due to incomplete knowledge or uncertain environment which implies that there exist various external sources and types of uncertainty. Fuzzy sets and fuzzy logic can be used to tackle uncertainty inherent in actual scheduling problems. Here, we use triangular fuzzy membership function to represents the uncertainty involved in processing of jobs.

Further, the system characteristics are described by membership function; it preserves the fuzziness of input information. However, the designer would prefer one crisp value for one of the system characteristics rather than fuzzy set. In order to overcome this problem we defuzzify the fuzzy values of system characteristic by

using the Yager's (1981) approximation formula. For a triangular fuzzy number $\tilde{A} = (a_1, a_2, a_3)$;

$$\text{crisp}(\tilde{A}) = \text{Average High Ranking of } A = h(A) = \frac{3a_2 + a_3 - a_1}{3}.$$

4. THEOREMS

The following theorems have been established to optimize the bicriteria scheduling on parallel machines involving maximum tardiness and weighted flow time

4.1. Theorem: A sequence S of jobs following early due date (EDD) order is an optimal sequence with maximum tardiness (T_{max}).

i.e. when jobs are processed on any of available parallel machines by early due date rule, the corresponding sequence of job processing is optimal with respect to maximum tardiness as given by Sharma et al (2013a).

4.2. Theorem: A sequence S of jobs following weighted smallest processing time (WSPT) rule, in which the jobs are placed to the earliest available location on the machines in WSPT order, minimizes the weighted flow time.

Proof: Let, if possible sequence S obtained by using the WSPT rule (i.e. by arranging the jobs in decreasing order of their weights; break the ties (if any) arbitrary) is not optimal. Let there exist a better sequence of jobs S' (say) in which adjacent jobs *i* and *j* are interchanged.

Let $C_i(S)$ and $C_j(S)$ be the completion times of jobs *i* and *j* in schedule S respectively. Similarly, let $C'_i(S)$ and $C'_j(S)$ be the completion times of jobs *i* and *j* in schedule S'.

Therefore, we can have: For sequence S: we have:

$$C_i(S) = t_a + 1, C_j(S) = t_a + 2$$

For sequence S': we have:

$$C_i(S') = t_a + 2, C_j(S') = t_a + 1$$

Next, the WTF contribution of these jobs for the sequence S is:

$$\begin{aligned} W(S) &= w_i \times C_i(S) + w_j \times C_j(S) \\ &= w_i \times (t_a + 1) + w_j \times (t_a + 2) \\ &= w_i \times t_a + w_i + w_j \times t_a + 2w_j = (w_i + w_j)t_a + w_i + 2w_j \end{aligned} \quad \text{---- (5)}$$

Similarly, the WTF contribution of these jobs for the sequence S' is

$$\begin{aligned} W(S') &= w_i \times C_i(S') + w_j \times C_j(S') = w_i \times (t_a + 2) + w_j \times (t_a + 1) = w_i t_a + 2w_i + w_j t_a + w_j \\ &= (w_i + w_j)t_a + 2w_i + w_j \end{aligned} \quad \text{---- (6)}$$

Since, the jobs *i* and *j* are placed by WSPT rule. Therefore, we have $w_i \geq w_j$. Hence, from results (5) and (6), we have: $W(S') \geq W(S)$. Therefore, WTF for the sequence S' is more as compared to the sequence S.

Hence, the sequence S following the Weighted Shortest Processing Time (WSPT) rule minimizes the Weighted Flow Time (WTF).

4.3. Theorem: A set of jobs initially arranged by Early Due Date order then a late job need to be considered for being exchanged only with another late job or a job having the same due date improves the value of secondary criteria of weighted flow time given the primary criterion of minimum total tardiness.

Proof: Let us pick any two jobs *i* and *j* from the EDD schedule.

Let job *j* be the late job. The following cases may arise:

Case I: Job *i* is late and $d_i < d_j$

In this case, we have either $c_i = c_j$ or $c_i < c_j$.

If $c_i = c_j$, then the switching of these jobs will not improve the solution.

If $c_i < c_j$, then the tardiness T and T' before and after the exchange are

$$T_i = \max(0, c_i - d_i) + c_j - d_j, T'_i = \max(0, c_i - d_j) + c_j - d_i$$

In case if $c_i > d_j$, then switching job *i* and job *j* will worsen the primary criterion.

In case if $c_i < d_j$, then switching job *i* and job *j* does not change the total tardiness and weighted flow time. The only case in which the primary criterion is not violated and weighted flow time improves is, if $c_i = d_j$.

Case II: If job *i* is not late and $d_i < d_j$

In this case, the total tardiness before and after switching job *i* and *j* is $T = c_j - d_j, T' = c_i - d_j$. Here, we have $T < T'$.

Hence, the primary criterion of total tardiness is violated.

Case III: If job *i* is not late job $d_i > d_j$

In this case, the total tardiness before and after switching job *i* and job *j* is $T = c_j - d_j, T' = c_i - d_j$. Here, we have $T < T'$.

Hence, the primary criterion of total tardiness is again violated.

Case IV: If job *i* is late and $d_i > d_j$

In this case, we get the similar result as we get in case I, discussed above.

Hence, we have shown that a switching among any two jobs will worsen the EDD schedule except that made under the exchange condition $c_i = d_j$ as stated in the algorithm. Hence, a set of jobs initially arranged in EDD order, a late job needs to be considered for being exchange only with another late job or a job having the same due date to potentially improve the value of a secondary criterion, given the primary criterion of minimizing total tardiness.

5. ALGORITHM

The following algorithm is proposed to find the optimal sequence for bi-criteria problem WFT/Total Tardiness:

Step 1: Find the crisp values of the fuzzy processing time of various jobs on different machines using Yager (1981) approximation formula.

Step 2: Arrange all the jobs on the available parallel machines in an early due date (EDD) order. If there is a tie then use Weightage Shortest Processing Time (WSPT) to break the tie.

Step 3: Let C be a set of jobs that cannot be switched and L be a set of all late jobs.

Initially, $C = \{\emptyset\}$.

Step 4: Consider first late job $i \notin C$. If none exist then go to the step 6; else go step 5.

Step 5: Check all the late jobs after job i . If there exist $j \in L, j \neq i$ such that $c_i \geq d_j$ and $w_i < w_j$, So we exchange the job i with job j which has maximum weight amongst all jobs satisfying these conditions; update L, if necessary else set $C = C + \{i\}$ and go to Step 4.

Step 6: Consider the first non late job $i \notin C$. If none exist then exit else go to step 7.

Step 7: Check all the early jobs after job i . If there exist a non late job j after job i for which $c_i \leq d_j$ and $c_j \leq d_i$ and $w_i < w_j$; exchange the job i with j which has maximum weight amongst all jobs satisfying the above conditions else set $C = C + \{i\}$ and go to Step 6.

6. NUMERICAL ILLUSTRATION

The following numerical illustrations are carried out to test the efficiency of algorithm proposed to optimize the bicriteria WFT/Total Tardiness on parallel machines in fuzzy environment.

6.1. Consider an example of jobs with processing time in fuzzy environment, due date and jobs Weight given in table 1 to be scheduled on three parallel machines in a flowshop. The objective is to obtain a Sequence of the jobs processing optimising the bicriteria taken as WTF/Total Tardiness.

Table 1. Jobs with fuzzy processing time.

Jobs	Processing Time	Due Date	Weight (w_i)
1	(8,9,10)	20/3	4
2	(5,6,7)	29/3	6
3	(9,10,11)	32/3	3
4	(7,8,9)	26/3	5
5	(5,6,7)	25/3	1
6	(10,11,12)	35/3	2

Solution: The crisp values for fuzzy processing time of given jobs is as follow

Table 2. Jobs with crisp values for processing time.

Jobs	Processing Time	Due Date	Weight(w_i)
1	29/3	20/3	4
2	20/3	29/3	6
3	32/3	32/3	3
4	26/3	26/3	5
5	20/3	25/3	1
6	35/3	35/3	2

On arranging the jobs in EDD order on the parallel available machines M_1, M_2 and M_3 , We have

Table 3. Jobs scheduling following EDD order.

Jobs	M_1	M_2	M_3	w_i	d_i	T_i
1	0 – 29/3			4	20/3	9/3
5		0 – 20/3		1	25/3	-
4			0 - 26/3	5	26/3	-
2		20/3 – 40/3		6	29/3	11/3
3			26/3 -58/3	3	32/3	26/3
6	29/3 – 64/3			2	35/3	29/3

Therefore, Total Tardiness = $\frac{9+11+26+29}{3} = \frac{75}{3}$ units and weighted flow time

$$\text{Weighted Flow Time} = \frac{\sum_{i=1}^n w_i c_i}{\sum_{i=1}^n w_i} = \frac{\frac{29}{3} \times 4 + \frac{20}{3} \times 1 + \frac{26}{3} \times 5 + \frac{40}{3} \times 6 + \frac{58}{3} \times 3 + \frac{64}{3} \times 2}{6 + 5 + 4 + 3 + 2 + 1} = 15.34 \text{ units.}$$

Therefore, set of late jobs = $L = \{1,2,3,6\}$ and set of jobs that cannot be switched $C = \emptyset$

On considering 1st late job $i = 1 \in L$ and $1 \notin C$ there is a late job $j = 2 \in L, j \neq i$ after job i such that $c_i \geq d_j$ and $w_i < w_j$. Therefore on interchanging the job i with job j , the job schedule becomes

Table 4. Reduced Job scheduling table.

Jobs	M_1	M_2	M_3	w_i	d_i	T_i
2	0 – 20/3			6	29/3	-
5		0 – 20/3		1	25/3	-
4			0 - 26/3	5	26/3	-
1	20/3 – 49/3			4	20/3	29/3
3		20/3 – 52/3		3	32/3	20/3
6			26/3 - 61/3	2	35/3	26/3

Therefore, Total Tardiness = 75/3 units and weighted flow time

$$WFT = \frac{\sum_{i=1}^n w_i c_i}{\sum_{i=1}^n w_i} = \frac{\frac{20}{3} \times 6 + \frac{20}{3} \times 1 + \frac{26}{3} \times 5 + \frac{49}{3} \times 4 + \frac{52}{3} \times 3 + \frac{61}{3} \times 2}{6 + 5 + 4 + 3 + 2 + 1} = 11.8 \text{ units.}$$

Therefore, set of late jobs = $L = \{1,3,6\}$ and set of jobs that cannot be switched $C = \emptyset$

On considering 1st late job $i = 1 \in L$ and $1 \notin C$, there is no late job after job i satisfying $c_i \geq d_j$ and $w_i < w_j$, therefore $C = C + \{i\} = \{1\}$.

Now, on considering the next late job $i = 3 \in L$ and $3 \notin C$, there is no late job after job i satisfying $c_i \geq d_j$ and $w_i < w_j$, therefore $C = C + \{i\} = \{1,3\}$.

Now, on considering the next late job $i = 6 \in L$ and $6 \notin C$, there is no late job after job i .

So set $C = C + \{i\} = \{1,3,6\}$.

Now, there is no late job $i \notin C$, so we pick the first non late job $i = 2 \notin C$.

Next, we observe that there is no early job j after job i in the schedule for which $c_i \leq d_j$ and $c_j \leq d_i$ and $w_i < w_j$; so set $C = C + \{i\} = \{1,2,3,6\}$

Now, on consider the next non late job $i = 5 \notin C$, there is no early job j after job i in the schedule for which $c_i \leq d_j$ and $c_j \leq d_i$ and $w_i < w_j$, so set $C = C + \{i\} = \{1,2,3,5,6\}$

Similarly, on consider the next non late job $i = 4$, there is no early job j after job i in the schedule.

Therefore $C = C + \{i\} = \{1,2,3,4,5,6\}$

Hence, the optimal sequence of jobs processing is 2 – 5 – 4 – 1 – 3 – 6 with minimum WFT as 11.8 units and Total tardiness as 75/3 units.

7. DISCUSSION

7.1. The proposed algorithm optimizes the bi-criteria problem of minimizing the weighted flow time for a minimum value of total tardiness.

Proof: The result can be verified by considering the following two cases:

Case 1: When jobs i and j are late jobs, i.e. i and $j \in L$

This case corresponds to step 5 of the algorithm. We know that if jobs are initially arranged in early due date order, a late job need to be considered for being exchanged only with another late job in order to potentially improve the value of secondary criterion of weighted flow time, given the primary criterion of minimum total tardiness.

If the conditions $c_i \geq d_j$ and $w_i < w_j$ for $i, j \in L$, then job j appearing after job i in the schedule violate the primary criterion of minimum total tardiness. Hence, jobs i and j must be exchanged in order to optimize the secondary criterion of weighted flow time for a given minimum value of primary criterion of total tardiness.

Case 2: When jobs are early jobs

This case corresponds to step 7 of the algorithm. Since, early jobs remain early jobs even when they are exchanged. Therefore, if there exist a non late job j after an early job i for which $c_i \leq d_j$ and $c_j \leq d_i$ and $w_i < w_j$ then interchange the job i with job j which has maximum weight amongst all jobs satisfying the above conditions otherwise, the secondary criterion of weighted flow time will remain optimized with minimum total tardiness.

7.2. If the problem of single criteria, Total Tardiness, is NP-hard, the scheduling problem on parallel machines optimising the bi-objective function WFT / Total Tardiness will also be NP-hard.

Solution: We shall prove the result by the method of contradiction:

Let if possible the bi-objective function WFT / Total Tardiness is not NP-hard. Therefore, there must exist a polynomial algorithm which can solve the problem of optimising the bi-objective function WFT / Total Tardiness on parallel processing machines.

This implies that single criteria of Total Tardiness can be optimized in polynomial time; i.e. Total Tardiness is not NP-hard. This is a contradiction as Total Tardiness is NP-hard.

Hence, the scheduling problem optimising the bi-objective function WFT / Total Tardiness on parallel processing machines will also be NP-hard.

8. CONCLUSION

In this paper a heuristic algorithm to optimize the bicriteria scheduling problem involving total tardiness and weighted flow time on parallel machines is developed. In real life situation the processing time of jobs may vary due to lack of complete knowledge, uncertainty and vagueness. The concept of fuzzy processing time is introduced in processing of jobs to handle under these uncertainties. The optimal sequence of jobs processing for the problem discussed above is 2 – 5 – 4 – 1 – 3 – 6 with minimum WFT as 11.8 units and Total tardiness as 75/3 units. The study may further be extended by using trapezoidal fuzzy numbers, by generalizing the number of machines, by introducing the concepts of non availability constraints and machines processing the jobs with different speeds.

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